On the Semantics of Mondel Specifications Using a Concurrent Object Petri Net-Based Approach

Nasreddine Aoumeur

ITI, FIN, Otto-von-Guericke-Universität Magdeburg
Postfach 4120, 39016 Magdeburg, Germany
E-mail: {aoumeur}@iti.cs.uni-magdeburg.de
Phone: ++49/391/67-{18066|12994|18800} Fax: ++49/391/67-12020

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Chapter 1

Abstract

The purpose of the present paper is twofolds. Firstly, we present our first results towards a tailored specification and validation object-oriented (OO) model for distributed systems. Referred to as CO-Nets, the model is a variant of object Petri nets characterized by: a complete integration of object-oriented concepts and constructions into an appropriate form of algebraic Petri nets named ECATNets; two communication patterns for intra- and inter-objects interaction that enhance modularity and concurrency without violating the encapsulated aspects of each class and last but not least, the interpretation of the behaviour of the constructed model into rewriting logic which allows some validation by rapid-prototyping using rewrite techniques. Secondly, as a significant case study, we assess the suitability of the CO-Nets approach by providing Mondel OO specification language with a formal semantics. Mondel language has been conceived for developing distributed applications with particularly: a state-oriented style of description, synchronous communication based on the rendez-vous mechanism and a formal semantics, based on coloured Petri nets, that we propose to improve using the CO-Nets approach.

Contents

Chapter 2

Introduction and motivation

The object orientation, with its powerful abstraction mechanisms, is nowadays a widely accepted paradigm dealing with different phases i.e. analysis, specification/validation, design and implementation in the development life cycle of large software in general and distributed ones particularly. Indeed, the object paradigm allows for: an intrinsic concurrency and distribution through its message passing mechanism and information hiding; a natural conceptualization of systems as a community of objects and messages, which is very close to the intuitive perception of distributed systems; an incremental building of complex systems through its powerful abstraction mechanisms (including object composition, inheritance and interaction).

These advantages and appropriateness have been fully exploited and confirmed by a great number of OO approaches, languages and formalisms for developing large systems in general and distributed ones in particular. Among them we have especially: OOD [Boo91] and OMT [RBP$^+$91] as OO approaches and methodologies; TROLL [JSHS96], Maude [Mes93] and Mondel [BBE$^+$91] as OO specification/programming languages; and HOSA [GD93], rewriting logic [Mes92] and recently DTL [ECSD98] as (algebraic based) formalisms.

The Mondel OO specification language is particularly distinguished by the following features: object description as type instance with a persistent identity; multiple inheritance based on subtyping; concurrency with synchronous communication based on the rendez-vous mechanism and a state-oriented methodology for developing applications. Beside that, Mondel specification has received a formal semantics [BB91], using coloured Petri nets, that allows verification of some crucial properties of the specification like absence of deadlock, executability of operations, etc, on the basis of coloured Petri nets analysis properties [Jen92].

However, in spite of the strengths of the proposed Mondel formal semantics, we claim that a more appropriate semantics can be obtained by overcoming the following shortcomings, that present this formalization, both at the translating ideas level as well at the model level.

- In the modeling of tokens in places as object instances of the form $< Id, attributes, stack >$, the Id(object identifier) and $attributes$(object attributes names with their actual values) components are very intuitive and sound translation ideas; however, the *stack component* — storing values of actual operation parameters, local variables and intermediate computations— is some how artificial, which make any straightforward mechanization of the translation difficult. Moreover, the messages (i.e. operations on objects), that have to be independent entities (i.e. may be created, deleted, etc), are rather conceived as fixed part of the stack content, which avoids any form of concurrent processing.

- The separation between object acquaintances (i.e. attributes) and its 'controlled' internal states, and the modeling of each object internal state as a place seems also to be non-intuitive translation ideas. Indeed, as will be shown later in our approach, the perception of such internal states *just like* the other (stateless) attributes is more natural. On the other hand, by avoiding these (state) places, the resulting net is more simple, manageable and understandable.
- CPNets have been suitably used for modelling object instances as 'coloured' tokens and object behaviour as appropriate transitions. On the other hand, some Mondel specification properties have been verified using CPNets properties analysis. However, CPNets need for an appropriate (OO) extensions, like the one given by C. Lakos [Lak95] for example, for dealing with the *inheritance* and the persistence of *object identities* that are the main concepts of the OO paradigm in general, and Mondel specification particularly.

With the aim to cope with these shortcomings, the present paper proposes to translate Mondel specifications into the CO-Nets approach that we are currently developing. The CO-Nets OO approach is intended to modeling and rapid-prototyping distributed systems with more adequacy to information systems specification. As modeling framework, CO-Nets model is based on a judicious integration of OO concepts and constructions (including classes as modules, inheritance and interaction) into the ECATNets [BM92] algebraic Petri nets, which is mainly distinguished by its capability of separating between the enabling tokens and the destroyed ones and by its true concurrent semantics expressed into rewriting logic. For rapid-prototyping, CO-Nets behaviour is expressed into rewriting logic, and hence rapid-prototypes can be directly generated using rewriting techniques.

The structure of the remaining sections is as follows: in section 2, the main concepts of the ECATNets model are reviewed. The third section describes the CO-Nets approach, but without entering into technical details (a more formal presentation can be found in [AS99]). In section 3, we explain the basic features of the Mondel language using the vending machine example. In the main section, with the help of the vending machine example, the main ideas of translating Mondel into CO-Nets are detailed. Section 5 closes the paper by giving some concluding remarks.

For the rest of the paper we assume the reader is more or less familiar with algebraic Petri nets and rewriting techniques and logic. Good references for these topics are respectively [EM85], [JR91] and [Rei91] for the algebraic setting and the algebraic Petri nets, [DJ90] and [Mes92] for rewriting techniques and logic. Throughout the paper, we use for the algebraic description of different structures an OBJ notation [GWM$^+$92].

Chapter 3

ECATNets : An Overview

The ECATNets [BMB93] model has been conceived around three formalisms, the first two ones constitute a net/data model and serve for describing the structural aspects of the modeled system, while the third one allows for defining the semantics (i.e. the behaviour) of the system expressed into rewriting logic which will be first reviewed.

3.1 Rewriting logic

Rewriting logic, as a new paradigm for concurrent systems, have been introduced by J.Meseguer in[Mes92] by observing, first that *concurrent* rewriting is a natural process in terms rewriting and second the inadequacy of interpreting rewrite rules as (oriented) equations when dealing with non Platonic(i.e. reactive) systems. Then, in rewriting logic while rewrite rules have the usual form, they are rather interpreted as a *change* (i.e. becoming) in concurrent systems. To be more precise we present some definitions borrowed from[Mes92]:

• A *(labelled) rewrite theory* $\mathcal{R}$ is 4-tuple $\mathcal{R}= (\Sigma, E, L, R)$ where Σ is a ranked alphabet of functions symbols, E is a set of Σ-equations, L is a set called the set of *labels,* and R is a set of pairs $R \subseteq L \times (T_{\Sigma,E}(X)^2)^+$ whose first component is a label, and whose second component is a nonempty sequence of pairs of E-equivalence classes of terms, with $X = \{x1, .., xn\}$ a countably infinite set of variables. Elements of R are called *rewrite rules.* A rewrite rule $(r, [t], [t'])([u_1], [v_1]. . .([u_k][v_k])$ is denoted

$$r : [t] \Rightarrow [t'] \; if \; [u_1] \Rightarrow [v_1] \; \wedge \; \ldots \wedge \; [u_k] \Rightarrow [v_k]$$

• Given a rewrite theory $\mathcal{R}$, we say that $\mathcal{R}$ *entails* a sequent $r : [t] \Rightarrow [t']$ and write $\mathcal{R} \vdash [t] \Rightarrow [t'] \; iff \; [t] \Rightarrow [t']$ can be obtained by finite application of the following *rules of deduction*:

1. **Reflexivity** : For each [t] $\in T_{\Sigma,E}(X)$, $\frac{}{[t]\Rightarrow[t]}$
2. **Congruence** : For each f $\in \Sigma_n$, n $\in$ N, $\frac{[t_1]\Rightarrow[t'_1]. \;. \;.[t_n]\Rightarrow[t'_n]}{f(t_1,...,t_n)]\Rightarrow[f(t'_1,...,t'_n)]}$
3. **Replacement** : For each rule $r : [t(\bar{x})]$[1] $\Rightarrow$[t'$(\bar{x})$] $if \; [u_1(\bar{x})] \Rightarrow [v_1(\bar{x})] \; \wedge \; ..\wedge \; [u_k(\bar{x})] \Rightarrow [v_k(\bar{x})] \; in \; R$,

[1] $\bar{x}$ is just a simplified notation of $x1, ..., xn$

if $[w_1] \Rightarrow [w_1^{'}]...[w_n] \Rightarrow [w_n^{'}]$ then $\frac{[u_1(\bar{w}/\bar{x})]\Rightarrow[v_1(\bar{w}/\bar{x})]\wedge\ldots\wedge[u_k(\bar{w}/\bar{x})]\Rightarrow[v_k(\bar{w}/\bar{x})]}{[t(\bar{w}/\bar{x}]\Rightarrow[t(\bar{w}^{'}/\bar{x}]}$

4. **Transitivity** : $\frac{[t_1]\Rightarrow[t_2] \qquad [t_2]\Rightarrow[t_3]}{[t_1]\Rightarrow[t_3]}$

• Given a rewrite theory $\mathcal{R}= (\Sigma, E, L, R)$, a (Σ, E)-sequent $[t] \Rightarrow [t']$ is called : a *concurrent* $\mathcal{R}-$*rewrite*[2](or just a *rewrite*) iff it can be derived from $\mathcal{R}$ by finite application of the rules 1-4.

3.2 ECATNets : Structural aspects

As a usual Petri net, an ECATNet is a digraph with two kinds of vertices, called places and transitions, respectively represented by circles and boxes. However, tokens contained into ECATNets places are of *complex form.* More precisely, tokens into places[3] are (ground algebraic) terms following a given user signature [EM85]. Places are connected to transitions by directed (input) edges labelled of by two informations: The input Conditions(IC) and the Destroyed Tokens(DT) which are both multisets (i.e. set with possibility of repetition) of terms. Also, transitions are connected to places by directed (output) edges labelled of by a Created Tokens(CT) multiset. Finally, with each transition is associated a boolean expression called Transition Condition(TC). The generic ECATNets [BM92] model is depicted in figure 1(a). Note that, Destroyed Tokens(DT) are omitted when they are equal to Input Condition(IC).

The ECATNets operational behaviour is defined as follows: the firing conditions of transition t are : first, for all input places p to the transition t, $IC(p,t)$ are included in $M(p)$; second the transition condition of t is valued to true. After firing the transition t, we have simultaneously deletion of Destroyed Tokens from the corresponding (input) places and addition of the Created Tokens to the corresponding (output) place.

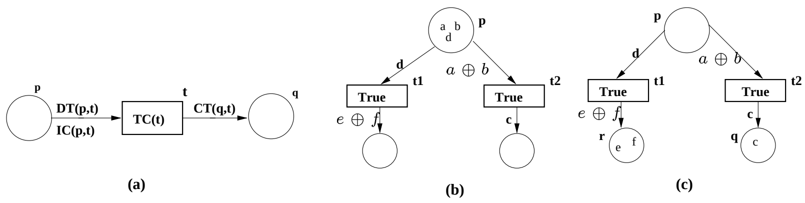

Figure 3.1: The Generic ECATNets and illustrative example

Example 3.2.1 *Assuming a, b, c, d, e and f as algebraic (constant) terms with respect to a given signature, and $\oplus$ as multiset union symbol on algebraic terms, figure 1(b) depicts a simplified ECATNet* □

[2]Noting that, other more restricted kinds of sequent are possible[Mes92]

[3]as usually, tokens into a given place p are denoted by M(p)

3.3 ECATNets : Semantical Aspect

The ECATNet semantics is expressed into rewriting logic, where the effect of each transition is captured by a rewriting rule in this logic. Depending on the relation between the Input Condition and the Destroyed Tokens, four cases have been distinguished[BMB93]. We recall the corresponding rewrite rules for these cases after giving some necessary notations:

- as described above, the content of each place is a multiset of closed terms w.r.t. a given signature. The usual union (assoc. and commu.) operation on this multiset is denoted by $\oplus$.
- to capture rewrite rules and ECATNets states (i.e. tokens distribution), another multisets form having as elements pairs of the form: (*place_name, multiset of terms*) was used. The union operation on this multiset is denoted by $\otimes$.
- given a transition t, let $\{p1, .., pn\}$ its input places and $\{q1, .., qm\}$ its output places.

case1: $IC(p,t) = DT(p,t)$, for each $p \in \{p1, ..., pn\}$. In this case the rule takes the form:
$t : (p1, IC(p1,t)) \otimes .. \otimes (pn, IC(pn,t)) \Rightarrow (q1, CT(q1,t)) \otimes .. \otimes (q_m, CT(q_m,t)\ if\ TC(t)$.

case2: $IC(p,t) \cap DT(p,t) = \emptyset$, for each $p \in \{\ p1, ..., pn\}$. In this case the rule takes the form:
$t : (p1, IC(p1,t)) \otimes (p1, DT(p1,t)) \otimes .. \otimes (pn, IC(pn,t)) \otimes (pn, DT(pn,t) \Rightarrow$
$(p1, IC(p1,t)) \otimes (q1, CT(q1,t)) \otimes .. \otimes (q_m, IC(q_m,t)) \otimes (q_m, CT(q_m,t)\ if\ TC(t)$.

case3: $IC(p,t) \cap DT(p,t) \neq \emptyset$, for at least one $p \in \{p1, .., pn\}$. It has been shown in [BM92] that this case could be brought to the two already cases.

case4: The first three cases concern the CATNet version [BM92]: For the E(xtended) CATNets, two additional notations have been introduced [BMB93]. Both notations concern the Input Condition. In the first one IC takes the form: $IC(p,t) =^{\sim}$*multiset of terms*. Where $^{\sim}$ is a new symbol. The enabling condition in this case means that $IC(p,t)$ has not to be included in $M(p)$. Using the second notation we may have $IC(p,t) = \emptyset$ which simply means that the associated transition may be fired only when its input places are empty. The associated rewrite rules are detailed in [BMB93].

Example 3.3.1 *For the above simple example, by applying the first case, we obtain the following rules for the transitions* $t1$ *and* $t2$. $t1 : (p,d) \rightarrow (r, e \oplus f)$ $t2 : (p, a \oplus b) \rightarrow (q,c)$. *The intial state, represented in figure* $2(b)$*, is equivalent to:* $(p, a \oplus b \oplus d) \otimes (q, \emptyset) \otimes (r\emptyset)$. *Deduction of the next state depicted in figure 1(c) may be obtained by the following rewrite sequent:* $(p, a \oplus b \oplus d) \otimes (q, \emptyset) \otimes (r, \emptyset) \Rightarrow (p, a \oplus b) \otimes (p, d) \otimes (q, \emptyset) \otimes (r, \emptyset)$ *"By applying the decomposition rules [BM92]"* $\Rightarrow (q,c) \otimes (r, e \oplus f) \otimes (q, \emptyset) \otimes (r, \emptyset)$ *"By applying concurrently the rewrite rules associated with* $t1$ *and* $t2$*"* $\Rightarrow (q,c) \otimes (r, e \oplus f)$ *"by applying concurrently the commutativity and the structural axioms of identity [BM92]"* .

Chapter 4

The CO-Nets Approach

As mentioned above, the CO-Nets is an OO approach for developing distributed systems. In this sense, CO-Nets aim to cover the important phases in developing such systems, particularly the formal specification and the verification/validation phases. In what follows, we present the main aspects of this approach by putting more emphasis on the modelling phase. Firstly, we present how structural aspects of a given system are captured using suitable OO signature, we then deal with the representation of objects and classes into the CO-Nets model. More complex systems are modeled next, using more advanced abstraction mechanisms, especially the specialization and the interaction between components.

4.1 The CO-Nets: Structural Aspects

The first intuitive basic ideas for representing an OO system (i.e. a community of objects) into the ECATNets framework is to regard object states and message instances (i.e. method invocation) as algebraic terms— with adding logical identities to objects states and ensuring their uniqueness. Second, represent such message and state terms into associated places and their effect— as a result of interaction of messages with the objects to which they are sent— by corresponding transitions.

However, for the modelling of more complex systems as interacting components, we need further powerful mechanisms that allow for capturing OO constructions like: specialization, object composition and particularly the interaction between differents classes. For this aim, on the one hand, we propose to make clear distinction between local properties (i.e. objects attributes) that have to be hidden from the outside of the object and external ones that can be observed (and may be modified) by other objects classes. On the other hand, we propose as well to distinguish between the internal messages that allow for evolving the object states over the time and the external messages used for interacting different classes.

More precisely, each object state will be regarded as a term —with respect to the OO signature below— of the form $< I|atr_1 : val_1, ..., atr_k : val_k, atbs_1 : val'_1, ..., atbs_s : val'_s >$[1]; where I is an observed object identity taking its values from an appropriate abstract data type typed by OId; $atr_1, ..atr_k$ are the *local* attributes identifiers component that have as actual values respectively $val_1, .., val_k$. The observed part of an object state is identified by

[1]This structure is inspired by the Maude language [Mes93] but with an explicit distinction between hidden and observed —from the outside— attributes.

$atbs_1, .., atbs_s$ with their associated actual values $val'_1, .., val'_s$. Also, we assume that all the attributes identifiers(local or observed) range their value over a suitable sort denoted AID, and their associated values are ranged over the sort $Value$ with $OId < Value$ (i.e. OId as subsort of $Value$) to allow object valued attributes. Also, in order to have more flexibility of this object state and allow by the way to exhibit intra-object concurrency, we introduce an appropriate operator that we denote by $\oplus$ for splitting (resp. recombining), if need be, this state into part (resp. its part). This splitting/recombining operation, reflected by the axiom in the object-state signature as depicted below, is particularly important for the respect of the encapsulation property when interacting different classes.

Messages are viewed as operations with at least one of their arguments is of OId sort — in this sense a message should involve at least one object as sender or receiver. Each message generator ms_i will be typed by a sort denoted Ms_i. Moreover, as mentioned above, we distinguish between local and external messages. The local messages to a given class Cl have to include at least the two usual messages: the message destinated for creating a new object state and the message for the deletion of existing object that we denote respectively by Ad_{Cl} and Dl_{Cl}.

All for all, following these informal description and some ideas from [Mes93], the formal description of the object states as well as the class structures, using an OBJ [GWM+92] notation, is depicted in figure 4.1.

On the basis of the class description, we define informally the associated CO-Net structure as follows:

- The places of the CO-Net are precisely defined by associating with each message generator one place that we called message place. Therefore, each message place have to contain message instances, of a specific form, sent to the objects (and not yet performed). In addition to these message places, we associate with the object sort one (object) place that have to contain the current object states of this class.

- CO-Net transitions reflect the effect of messages on the object states to which they are sent. Also, we make distinction between local transitions that reflect the object states evolution and external transitions modeling the interaction between different classes. The conditions to be fulfilled for each kind of these transitions forms is given the subsection below.

- Conditions may be associated with transitions. They involve attribute and/or message parameters variables.

4.2 CO-Nets : Semantical Aspects

After giving how CO-Nets templates, as description of classes, are constructed, we focus herein on the behavioural aspects of such classes. That is, how to construct a *coherent* object society as a community of object states and message instances, and how such a society evolves only into a *permissible* society. By coherence we mean the respecting of the system structure and the uniqueness of objects identities. And by permission we mainly understand the respecting of the encapsulation property.

```
obj Object-State is
  sort AId .
  subsort OId < Value .
  subsort Attribute < Attributes .
  subsort Id-Attributes < Object .
  subsort Local-attributes External-attributes < Id-Attributes .
  protecting Value OId AId .
  op _:_ :  AId Value → Attribute .
  op _,_ :  Attribute Attributes → Attributes [associ. commu. Id:nil] .
  op ⟨_|_⟩ :  OId Attributes → Id-Attributes .
  op _⊕_ :  Id-Attributes Id-Attributes → Id-Attributes [associ. commu. Id:nil] .
  vars Attr:  Attribute ; Attrs$_1$, Attrs$_2$:  Attributes ; I:OId .
  eq1 $\langle I|attrs_1\rangle \oplus \langle I|attrs_2\rangle = \langle I|attrs_1, attrs_2\rangle$ .
  eq2 $\langle I|nil\rangle = I$
endo.
```

```
obj Class-Structure is
  protecting Object-state, s-atr$_1$,...,s-atr$_n$, s-arg$_{l1,1}$,.., s-arg$_{l1,l1}$,
             ...,s-arg$_{i1,1}$,...,s-arg$_{i1,i1}$ ...
  subsort Id.obj < OId .
  subsort Mes$_{l1}$, Mes$_{l2}$,...,Mes$_{ll}$ < Local_Messages .
  subsort Mes$_{e1}$, Mes$_{e2}$,...,Mes$_{ee}$ < Exported_Messages .
  subsort Mes$_{i1}$, Mes$_{i2}$,...,Mes$_{ii}$ < Imported_Messages .
  sort Id.obj, Mes$_{l1}$, .   .   .   ,Mes$_{ip}$
  (* local attributes *)
    op $\langle \_|atr_1 : \_, ..., atr_k : \_\rangle$ :  Id.obj s-atr$_1$ ...s-atr$_k$
        → Local-Attributes.
  (* observed attributes *)
    op $\langle \_|atrbs_1 :, ..., atrbs_{k'} : \_\rangle$ :  Id.obj s-atbs$_1$ ...s-atbs$_{k'}$
        → External-Attributes.
  (* local messages *)
    op ms$_{l1}$:  s-arg$_{l1,1}$ ...s-arg$_{l1,l1}$ → Mes$_{l1}$ .  ...
  (* export messages *)
    op ms$_{e1}$:  s-arg$_{e1,1}$ ...s-arg$_{e1,e1}$ → Mes$_{e1}$ .  ...
  (* import messages *)
    op ms$_{i1}$:  s-arg$_{i1,1}$ ...s-arg$_{i1,i1}$ → Mes$_{ip}$ .  ...
endo.
```

Figure 4.1: The Template Object Signature of Object Systems

4.2.1 Object creation and deletion

For ensuring the uniqueness of objects identities in a given class that we denote by Cl, we propose the following conceptualization:

1. Add to the associate (marked) CO-Net modeling such class, a new place of sort $Id.obj$ and denoted by $Id.Cl$ containing *actual objects identifiers* of objects of the place obj.

2. Objects creation is made through the net depicted in the left hand side of figure 4.2. Where the notation $\sim$ captures exactly the intended behaviour (i.e. the identifier Id should not already be in the place $Id.Cl$). After firing this transition, there is an addition of this new identifier to the place $Id.Cl$ and a creation of a new object, $< Id|atr_1 : in_1, ..., atr_k : in_k >$, with $in_1, ..., in_k$ as optional initial attributes values.

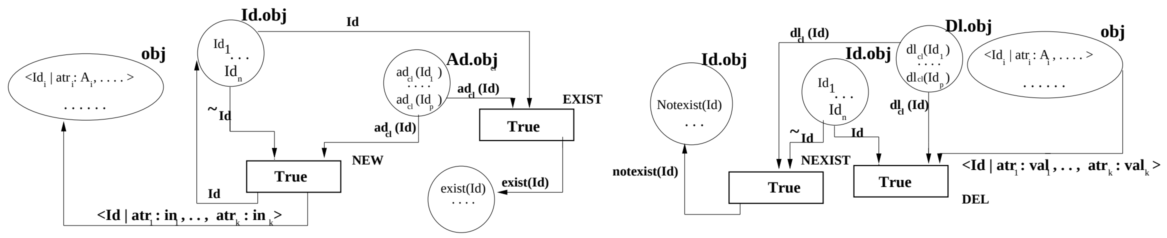

Figure 4.2: Object Creation and Deletion Using OB-ECATNets

This objects creation model is very general, and can be adapted to specific cases. For example, when restriction conditions should be associated with initial values, the signature of the creation message can be modified to: ad_{Cl} : $Id.obj\ s\_atr_1\ ..\ s\_atr_k \rightarrow ad_{Cl}$ which allow to associate conditions with the transition *NEW*.

As depicted in the right hand of figure 4.2, the objects deletion is modelled following the same reasoning.

4.2.2 Evolution of Object States in Classes

For the evolution of object states in a given class, we propose a general pattern that have to be respected in order to ensure the encapsulation property—in the sense that no object states of other classes participate in the communication — as well as the preservation of the object identity uniqueness. Following such guidelines and in order to exhibit a maximal concurrency, this evolution schema is depicted in Figure 4.3, and it can be intuitively explained as follows: The contact of the only relevant parts[2] of some object states of a given Cl, —namely $< I_1|attrs_1 >$[3] ;..; $< I_k|attrs_k >$— with some messages $ms_{i1},..,ms_{ip},ms_{j1},..,ms_{jq}$—declared as *local or imported* in this class— and under some conditions on the invoked attributes and message parameters results results in the following effects:

- The messages $ms_{i1},..,ms_{ip},ms_{j1},..,ms_{iq}$ vanish;
- The state change of some (parts of) object states participating in the communication, namely $I_{s1},..,I_{st}$. Such change is symbolized by $attrs'_{s1},..,attrs'_{st}$ instead of $attrs_{s1},..,attrs_{st}$.
- Deletion of the some objects by sending explicitly delete messages for such objects.
- New messages are sent to objects of the class Cl , namely $ms'_{h1},..,ms'_{hr},ms'_{j1},..,ms'_{jq}$.

4.2.3 Rewriting rules gouverning the CO-Nets behaviour

In the same spirit of the ECATNets behaviour, each CO-Net transition is captured by an appropriate rewriting rule interpreted into rewrite logic. Following the communication

[2]Such selection is possible due to the 'splitting/combination' axiom that have to be performed in front of each evolution depending on the invoked rewrite rule.

[3]$attrs_i$ is simplified notation of $atr_{i1} : val_{i1},..,atr_{ik} : val_{ik}$.

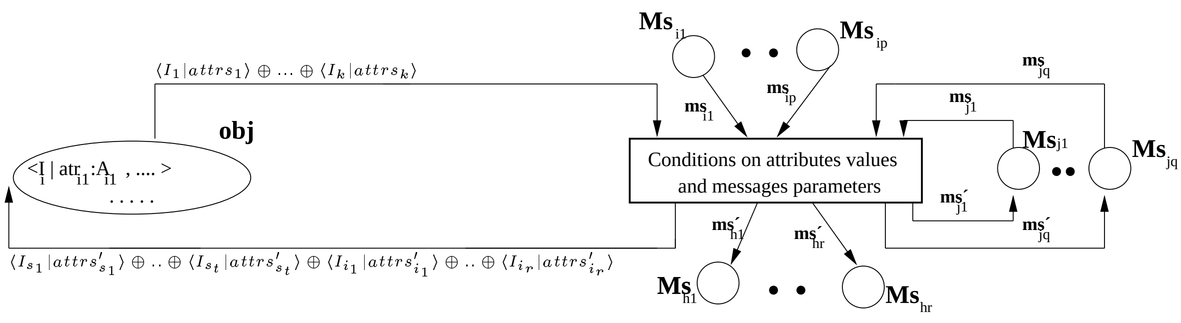
$\langle I_1|attrs_1\rangle \oplus ... \oplus \langle I_k|attrs_k\rangle$
obj
$<I_i | atr_{i1}:A_{i1}, ... >$
.....
$\langle I_{s_1}|attrs'_{s_1}\rangle \oplus .. \oplus \langle I_{s_t}|attrs'_{s_t}\rangle \oplus \langle I_{i_1}|attrs'_{i_1}\rangle \oplus .. \oplus \langle I_{i_r}|attrs'_{i_r}\rangle$
Ms_{i1}
Ms_{ip}
ms_{i1}
ms_{ip}
Conditions on attributes values and messages parameters
ms_{jq}
ms_{j1}
Ms_{j1}
Ms_{jq}
ms'_{j1}
ms'_{jq}
ms'_{h1}
ms'_{hr}
Ms_{h1}
Ms_{hr}

Figure 4.3: The Intra-Class CO-Nets Interaction Model.

pattern in figure 4.3, the general form of the rewrite rule associated with this intra-class interaction model takes the following form:

$$\begin{gathered}
T : (Ms_{i1}, ms_{i1}) \otimes .. \otimes (Ms_{ip}, ms_{ip}) \otimes (Ms_{j1}, ms_{j1}) \otimes .. \otimes (Ms_{jq}, ms_{jq}) \otimes \\
(obj, < I_1|attrs_1 > \oplus ... \oplus < I_k|attrs_k >) \Rightarrow (Ms_{h1}, ms'_{h1}) \otimes ... \otimes (Ms'_{hr}, ms'_{hr}) \otimes \\
(Ms_{j1}, ms'_{j1}) \otimes .. \otimes (Ms'_{jq}, ms'_{jq}) \otimes (obj, < I_{s1}|attrs'_{s1} > \oplus ... \oplus < I_{st}|attrs'_{st} >) \\
if\ Conditions \text{ and } M(Ad_{Cl}) = \emptyset \text{ and } M(Dl_{Cl}) = \emptyset\ ^4
\end{gathered}$$

4.3 CO-Nets : More Advanced Constructions

So far, we have presented only how the CO-Nets approach allows for conceiving independent classes. In what follows section, we give how more complex systems can be constructed using advanced abstraction mechanisms like the inheritance and the interaction between classes.

4.3.1 Inheritance in CO-Nets

In this paper, we restrict ourselves to the simple case of inheritance without overriding—the other forms of inheritance (i.e. multiple and with redefinition) may be found in [AS99]. Giving a (super)class modelled as a CO-Net, for constructing a subclass—with additional attributes and behaviour— that inherit the structure as well as the behaviour of its superclass and that it may exhibit new behaviour involving the additional attributes, we propose the following modeling. First, for allowing the inheritance of the structure and the behaviour of the superclass, we propose, on the one hand, to gather the object states of both classes (i.e. the super as well as the subclass objects) into the same object place (of the superclass). On the other hand, we add an *attribute variable* (i.e. variable of sort *Attributes* as declared in the object-state signature) to all arcs going to or from this place. Finally, the proper behaviour of the subclass is constructed in the usual way (i.e. by adding new places for each message and constructing the associated transitions reflecting the new behaviour that should take into consideration the additional attributes).

In a more clear way, if the IC, DT and CT associated with the arcs related to the object place contain terms of the form $< Id|atr_1 : V_1, ..., atr_k : V_k >$, we propose to transform them

[4]This condition requires that the creation and the deletion of objects have to performed at first. In other words, before performing this rewrite rules the marking in the Ad_{Cl} as well of the Dl_{Cl} places have to be empty.

into $< Id|atr_1 : V_1, ..., atr_k : val_k, ATS >$ where ATS can be instantiated either by *nil* whenever the objects of the superclass are involved by the associated transition or by the additional attributes(i.e. by $at\_ad_1 : Var_1, .., at\_ad_q : Var_q$)[5] whenever the subclass objects are involved. Noting that, $Var_1, .., Var_q$ should be declared as variables of *Attributes* sorts (see, the object-state signature). Semantically, each instantiation of the ATS variable is captured by a corresponding rewriting rule. Therefore, *two rewrite rules* are associated with each transition of the superclass.

4.3.2 Interaction in CO-Nets

Our conceptualization of the inter-class communication is inspired by the general communication schema proposed by J.Meseguer for concurrent OO systems [Mes93], but with taking into account the fact that intra-class objects evolution in each class is ensured by the model proposed in Figure 4.3. Therefore, as depicted in figure 4.4, the inter-class communication that we propose can be characterized as follows: The contact of some *imported or exported* messages namely $ms_{i1}, .., ms_{ip}, ms_{j1}, .., ms_{jq}$, that may belong to different classes denoted by $Cl_1, ..., Cl_m$, with exclusively the *external* parts of some objects states included in such classes, namely $< I_1|attrs_1 >$,..., $< I_k|attrs_k >$. Such contact, under some conditions that may involve objects attribute values and parameter messages, results in the following:

- The messages $ms_{i1}, .., ms_{ip}, ms_{j1}, .., ms_{iq}$ vanish;
- The state change of some (external parts of) object states participating in the communication, namely $I_{s1}, .., I_{st}$. Such change is symbolized by $attrs'_{s1}, .., attrs'_{st}$ instead of $attrs_{s1}, .., attrs_{st}$. The other objects components remain inchanged (i.e. there is no delection of parts of objects states).
- new external messages (that may involve delection/creaction ones) are sent to objects of different classes, namely $ms'_{h1}, .., ms'_{hr}, ms'_{j1}, .., ms'_{jq}$.

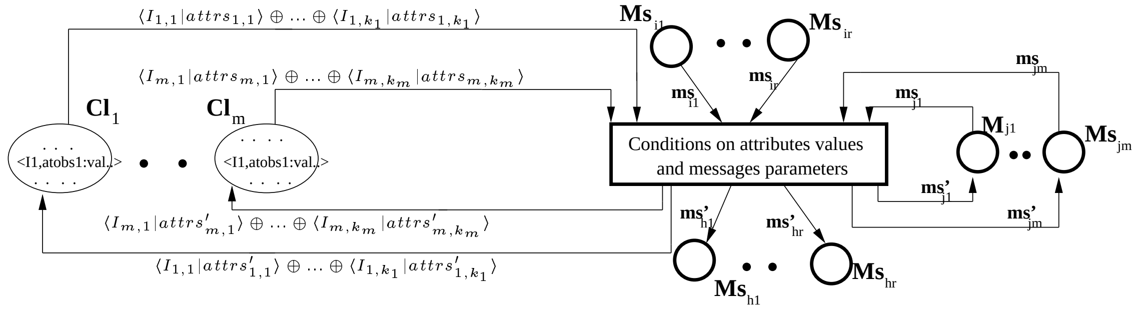
$\langle I_{1,1}|attrs_{1,1}\rangle \oplus ... \oplus \langle I_{1,k_1}|attrs_{1,k_1}\rangle$
$\langle I_{m,1}|attrs_{m,1}\rangle \oplus ... \oplus \langle I_{m,k_m}|attrs_{m,k_m}\rangle$
Cl_1
Cl_m
<I1,atobs1:val..>
<I1,atobs1:val..>
$\langle I_{m,1}|attrs'_{m,1}\rangle \oplus ... \oplus \langle I_{m,k_m}|attrs'_{m,k_m}\rangle$
$\langle I_{1,1}|attrs'_{1,1}\rangle \oplus ... \oplus \langle I_{1,k_1}|attrs'_{1,k_1}\rangle$
Ms_{i1}
Ms_{ir}
ms_{i1}
ms_{ir}
Conditions on attributes values and messages parameters
ms_{jm}
ms_{j1}
M_{j1}
Ms_{jm}
ms'_{j1}
ms'_{jm}
ms'_{h1}
ms'_{hr}
Ms_{h1}
Ms_{hr}

Figure 4.4: The Interaction General Model Between classes

[5]Where $at\_ad_q, ..., at\_ad_q$ are assumed to be the additional attributes of the subclass.

Chapter 5

Mondel Overview

The Mondel language has been designed particularly for modelling and specifying distributed applications into an object oriented setting. The main features of Mondel specification include: object persistence, multiple inheritance, rendez-vous communication between objects. An object, in Mondel, is an instance of a type (i.e., a class in other OO languages) that specifies the properties that characterize all its instances and the operations that can be accepted by the object following their specific behaviour. More precisely, each object type is described through three sections:

The Attributes Section: that may be empty, describes the attributes identifiers and their types (i.e., their sorts) that are often object-valued. Moreover, as mentioned above, one of the specific feature of Mondel specification is that some attributes, those modeling *internal objects states*, are implicit and their value depend on the actual procedure being executed (on the object).

The Operations Section: The definition of a given type include the declaration of operations (called methods in other OO languages) that may be invoked by objects of other types. Each operation has a fixed number of parameters(which may be of objects and/or predefined types).

The Behaviour Section: It describes, through some procedures, the execution of each operation . Noting that, the Mondel specification philosophy favors *communication* over computation. In other words, types behaviour is specified using a state-oriented style, where internal states are modeled as Mondel procedures. Such communication is made through a rendez-vous between the caller and the callee (*accept* for receiving and ! for sending), where the caller must *wait* until the callee has issued the corresponding *return* statement.

Example: Let us consider a vending machine [EDB93]which receives a coin and delivers candies to its users. The specification of the vending machine system consists of one module composed of two types: the type *Machine* and the type *User*, as described below using Mondel syntax. The relation between the types *User* and *Machine* is represented by the attribute *m* of type *Machine*, defined within the *User* type. The operations, *InsertCoin* and *PushAndGetCandy*, are specified within the operation clause (as shown in lines 2 and

4). Note that these operations are without parameters. The user is initially in a *Thinking* state, and when he decides to buy a candy he inserts a coin. After the coin has been accepted, the user enters the *GetCandy* state, ready to accept a coin. Once a coin is inserted, the machine accepts the coin and then it enters the *DeliverCandy* state. After the user has pushed the button of the machine, it delivers him/her a candy and it becomes *Ready* again.

```
unit VMsystem =
type Machine = object with
operation
  InserCoin ;
  PushAndGetCandy ;
behaviour
 Ready
where
 Procedure Ready =
  accept InsertCoin do
    return;
   end;
    DeliverCandy ;
  endproc Ready

  Procedure DeliverCandy=
  accept PushAndGetCandy do
   return;
  end;
    Ready
  endproc DeliverCandy
  endtype Machine
type User = object with
m : Machine ;
 behaviour;
  Thinking ;
 where;
 procedure Thinking =
  m!InserCoin ;
  GetCandy ;
 endproc Thinking ;

 procedure GetCandy
  m!InserCoin ;
  Thinking ;
 endproc GetCandy ;
 endtype User ;
{the vending machine system behaviour }
 behaviour
 define Amachine new(Machine) in
  eval new(User(Amachine)) ;
 end ;
 endunit VMsystem
```

Chapter 6

Translating Mondel Specifications into CO-Nets

In this section, we provide Mondel specifications with a clean semantics based on CO-Nets. For this aim, first, we present how object types including attributes, operations and behaviour are modeled using CO-Nets. Also, for the behaviour section, more emphasis will be done to the communication aspects; because we estimate that the modelling of the computational aspects, using the *define* statement, are more trivial. Second, we shortly comment how inheritance can be taken into account. Our translation ideas are illustrated using the vending machine example.

6.1 Translating Mondel Specifications into CO-Nets

According to the CO-Nets approach as presented in section 2, a Mondel specification may be translated into CO-Nets in straightforward manner. More precisely, with each type, we associate an CO-Net (i.e. a class) composed of:

- **One object states place:** containing object states having as attributes those declared explicitly in the attribute section plus *the implicit ones* modeling the states of the object, and which are usually initialized in the beginning of the behaviour section.

  Example 6.1.1 *For the machine type, we have no explicit attributes. However, each machine is characterized by its state (shortly, st) initialized by the value Ready that can become Deliver-Candy. Then, object instances of type machine, modeled by a corresponding place, should be declared as follows:*

  op $< \_|st : \_ >$: $Id.machime\ state \rightarrow Machine.$

  Where Id.machine stands for Machine identities sort, and state have two values {Rd, Dv} (i.e. Rd stands for Ready and Dv for Deliver-Candy). The same reasoning may be applied to the User type, where the machine m is an explicit attribute and the User state (shortly, su), that can be in Thinking(shortly Th) or Get-Candy (shortly, Gt), is an internal state.

- **Operations places:** For each declared operation, we associate a corresponding place. If an operation is declared as $opr(arg_1, .., arg_n)$, where $arg_1, ..., arg_n$ are types parameters, we associate to it the following OBJ declaration:

  $$\textbf{op}\ opr : arg_1\ ...\ arg_n\ Id.callee\ Id.caller \rightarrow Msg_i$$

  The addition of the identifier sorts of the caller and the callee is necessary. In fact, it allows us to know, when a message is accepted by an object (identified by *Id.callee*), which object is concerned by the associated *return* (i.e. which object receive the permission to continue its execution here the *Id.callee*).

  According to this, we should also add to each object type that send messages (i.e. containing the primitive '!') a place containing *return* messages declared as:

  $$\textbf{op}\ ret : Id.callee \rightarrow Msg_i$$

  Example 6.1.2 *Following the above translation ideas, the CO-Net associated with the vending machine is described below. For instance, the two machine operations, modeled by two corresponding places, are declared as follows. The InserCoint operation is declared as:*
  op $ins : Id.machine\ Id.User \rightarrow Msg_1$.
  Hence, ins(u,m), means that the User u inserts the button of the machine m. The PushandDeliverCandy operation is declared as:
  op $push : Id.machine\ Id.User \rightarrow Msg_2$.

  For the User type, we have no explicit operation and then no operations places. However, it contains (two occurrences of) a sent primitive, and therefore we have to conceive a 'return' place that store messages instance of the form:
  op $ret : Id.User \rightarrow Msg_3$.

- **A Transition is associated with each procedure:** According to the two CO-Nets communications patterns, each procedure describing state change of an object can be easily captured by a transition. Particularly, the rendez-vous communication is modeled as follows:

  1. An object changes its state by accepting some operation. Following the accept general form [BB91]: *accept OpName[***Provided** *PureExpr* **do** *stat* **end**. This behaviour is captured by a suitable transition that express the contact of *OpName* with the associated *object state*—both modeled as tokens into their associated places— under the (transition) condition *PureExpr*. The result of firing such transition (i.e. the created tokens) is described by the statement *stat*.

     Example 6.1.3 *If we take, for example, the ready procedure from the machine type:*
     8 **Procedure** *Ready* =
     9 **accept** *InsertCoin* **do**

```
10    return;
11   end;
12    DeliverCandy ;
13   endproc Ready
```

This procedure is captured by the transition **RD***, as depicted in the associated net below, that takes as input places the* **ins***(sertCoin) and the* **Mc***(i.e., machine) places and as output ones the* **Mc** *and the* **Ret***(i.e. return) places. More precisely, the contact of ins(u,m) with* $< m|st : Rd >$ *results in the change of the machine state to* $< m|st : Dv >$ *and a sent of a knowledge to the User u that the message is accepted, expressed by ret(u).*

2. The same reasoning may be applied to the statement of the form: *PureExpr ! OpName[PureExprList)].* More explicitly, here the caller *Id.caller* send the message *OpName(Id.Callee, PureExpr, PureExprList)* to the object place associated with *OpName.* This fact is captured by a transition that takes as input place the caller object place and as output place *OpName.*

Example 6.1.4 *The procedure Thinking described in the User type as:*

```
26   procedure Thinking =
27    m!InserCoin ;
28    GetCandy ;
29   endproc Thinking ;
```

is captured by the transition **Th***, having as input place the* **User** *place which send the message ins(u,m) to the machine. But according to the rendez-vous communication principle, the* **User** *must wait until the acceptance of such a message (i.e. the firing of the transition* **Ac-Is** *: accept-insert).*
In order to make this waiting state explicit, we propose to add to the User states two other values, namely $wt1$ *(i.e. wait1) and* $wt2$ *(wait2). Where,* $wt1$ *expresses that the User, after he send a* $ins(u, m)$ *message, is in a waiting state before he enters the GetCandy(Gc) state (i.e. before receiving the associated return expressed by the firing of the transition* **Ac-Is***). The* $wt2$ *expresses the waiting to enter the Thinking(Th) state as modeled by the transition* **Ac-Ps***(short for accept-push) (i.e. waiting for the acceptance of the message push(u,m)).*

Note that these intuitive suggestions should be taken into account in each similar case. In other words, for each send primitive occurrence (i.e. **!** operation) one should associate a waiting value. We can also, as done in [BB91], model all wait occurrence as one and associate with *each return* a corresponding place. However, in this case the number of places may become unmanageable.

All for all, the structural as well as the behavioural (i.e. the associated CO-Nets) aspects of the vending machine are described as follows:

• The machine and user states Specification:

obj MU-state **is**
 protecting string .

```
 sort st-user, st-mc, Id.machine, Id.User .
 subsort Id.machine, Id.User < String
 op Rd, Dv : → St-mc .
 op Th, Gc, wt1, wt2 : → St-user .
endo
```

• Specification of the Machine and User states

```
obj Machine is .
 protecting MU-state .
 sort Machine Msg1 Msg2 .
 op < _|st : _ > Id.machine St-mc → Machine .
 Imported Messages .
 op ins : Id.User Id.machine → Msg1 .
 op push : Id.User Id.machine → Msg2 .
 vars m:Id.machine
endo

obj User is
 protecting MU-state .
 sort User return .
 op < _|su : _, mc : _ > St-user Id.machine → User.
 Imported Messages .
 op ret : Id.User → return .
 vars u:Id.user
endo
```

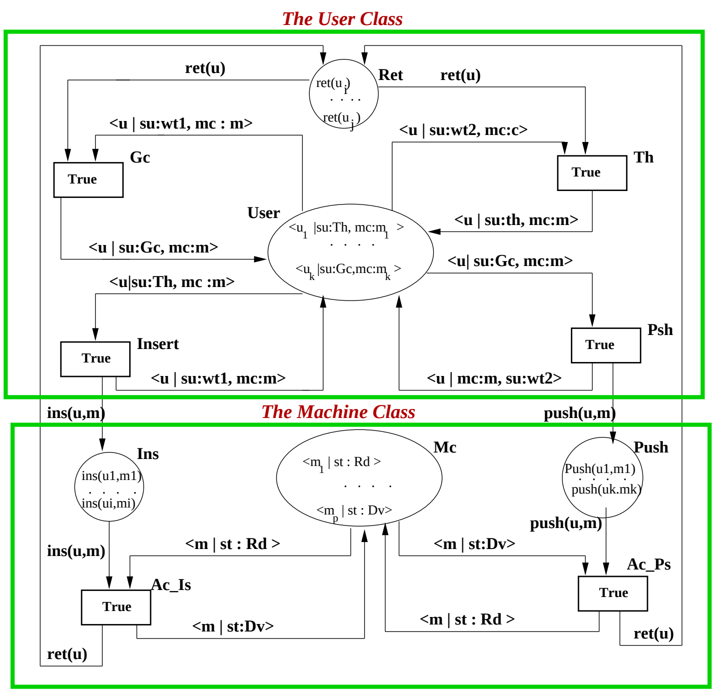

Figure 6.1: Modeling the VMachine Using CO-Nets

6.2 Translating Mondel subclasses into CO-Nets

Hereafter we explain how the proposed CO-Net inheritance allows for capturing the Mondel subtyping-based inheritance. For this aim, we propose to extend our running vending machine example. We assume that, beside the already specified machine, we may have a particular machine that before becoming ready it should be initialized. Moreover, we suppose that such initialization (operation) is possible only if the corresponding button is one-line. For capturing this informal description, we propose to add to the machine type new attribute named signal(shortly, *sg*), that have to be initialized by the value one(1), and a new state denoted by *init*. On the other hand, we assume that the user may send a message to initialize (shortly, *initialize*) such (special) machine. More precisely, this special vending machine can specified with Mondel as follows:

```
unit SVMsystem =
type SpMachine = Machine with
 sg : VAR[integer]
  operation ;
  initialize ;
newbehaviour
 init
where
 Procedure init =
  accept initialize do
    if sg = 1 then
    Ready
    endproc init
  endtype Machine

type User = object with
m : Machine ;
 newbehaviour;
  initial ;
 where;
 procedure initial =
  m!Initialize ;
  Thinking ;
 endproc initial ;
 endunit SpVMsystem
```

Following the CO-Nets inheritance principles, this vending machine extended behaviour can be easily captured as depicted in figure 5. Where the **Machine** place have to include now the usual machine states as well as the special machine states. The behaviour of the usual machine remains unchanged except, as required by the inheritance principle, the introduction of the ATS variable(that can take *nil* or $sg : Var$ as value) in the arcs related to the **Mc** place. For capturing the new behaviour in the machine subclass, a new place associated with the message *initialize* is constructed with its associated transition. Also, a new transition in the **User** class is added that reflect the sent of an initialization.

6.3 Rapid-Prototyping Mondel Specification using CO-Nets

As pointed out in the introduction, for proving crucial properties of a given specification using the CO-Nets approach rapid-prototypes can be generated using rewriting techniques and specifically the current implementation of the Maude language [MOM96]. Here we restrict ourselves to present the associated rewrite rules that reflect the CO-Nets behaviour associated with the vending machine system.

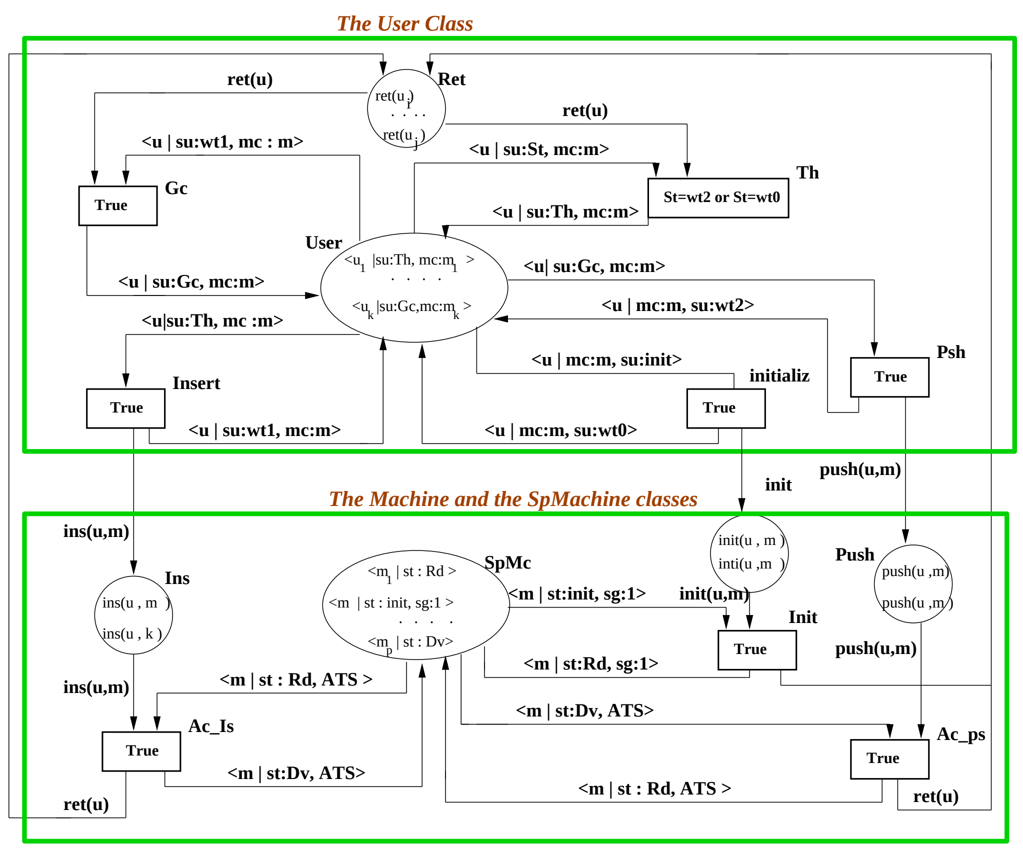

Figure 6.2: Modeling the SpVMachine Using CO-Nets

The User Class

Insert: $(User, < u|su : Th, mc : m >)$
$\Rightarrow (User, < u|su : wt1, mc : m >) \otimes (Ins, ins(u, m))$

Psh: $(User, < u|su : Gc, mc : m >)$
$\Rightarrow (User, < u|su : wt2, mc : m >) \otimes (Push, push(u, m))$

Gc: $(Ret, ret(u)) \otimes (User, < u|su : wt1, mc : m >)$
$\Rightarrow (User, < u|su : Gc, mc : m >)$

Th: $(Ret, ret(u)) \otimes (User, < u|su : wt2, mc : m >)$
$\Rightarrow (User, < u|su : Th, mc : m >)$

The Machine Class

Ac-Is: $(Ins, ins(u, m)) \otimes (Mc, < m|St : Rd >)$
$\Rightarrow (Mc, < m|st : Dv >) \otimes (Ret, ret(u))$

Ac-Ps: $(Push, push(u, m)) \otimes (Mc, < m|st : Dv >)$
$\Rightarrow (Mc, < m|st : Rd >) \otimes (Ret, ret(u))$

For the special Machine depicted in figure 5, the associated rewrite rules can be splitted in two classes: The rewrite rules that allows the inheritance of the behaviour of the machine and the rules that specify the new behaviour. The inherited part is captured by the following two rewrite rules

Ac-Is: $(Ins, ins(u, m)) \otimes (Mc, < m|St : Rd, sg : V >)$
$\Rightarrow (Mc, < m|st : Dv, sg : V >) \otimes (Ret, ret(u))$[1]

Ac-Ps: $(Push, push(u, m)) \otimes (Mc, < m|st : Dv, sg : V >)$
$\Rightarrow (Mc, < m|st : Rd, sg : V >) \otimes (Ret, ret(u))$

[1]Here the ATS attribute variable is instantiated by $sg : V$, where V is a variable that can take any variable of sg.

Chapter 7

Conclusion

In this paper, we proposed in a more intuitive way the CO-Nets approach as a formal model for concurrent OO systems. The approach is based on a complete integration of OO concepts and constructions into the ECATNets model, which is a form of high level Petri nets combining the strengths of Petri nets with those of abstract data type. Some key advantages of CO-Nets include: the modelling of simple and multiple inheritance in a straightforward way; the characterization of two communication patterns for intra-class as well as inter-class evolution promoting intra- and inter-objects concurrency and preserving the encapsulation property; the interpretation of the model into rewriting logic which allows the generation of rapid-prototypes using rewriting techniques and particularly the Maude language [Mes93].

As a significant case study for the assessment of the adequacy of the proposed model, for specifying and validating distributed systems in an object-oriented setting, we have shown how Mondel specifications can be easily and naturally translated into the CO-Nets framework. Moreover, due to the CO-Nets semantics, some of the Mondel properties can be verified either through graphical animation or by symbolic deduction using rewriting techniques. However, on the one hand, we have to investigate in a more detail which properties are particularly verifiable using these techniques. On the other hand, we plan to adapt the analysis techniques developed for coloured Petri nets and its object-oriented extensions developed in [Lak95].

Besides that, we are working for tailoring the CO-Net approach for the specification and validation of distributed information systems that present more difficulties at the structural level, which necessitate more advanced structuring mechanisms like the aggregation and the notion of role, as well as at the behaviour level, where notion of events composition have to be taken into account.

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